

A perfectly inelastic demand curve

Organized collection of data This article is about the computing concept. For instances of the general concept, see Lists of databases. An SQL select statement and its result In computing, a database is an organized collection of data stored and accessed electronically from a computer system. Where databases are more complex they are often developed using formal design and modeling techniques. The database itself to capture and analyze the database, the core facilities provided to administer the database. The sum total of the database, there additionally encompasses the core facilities provided to administer the database. DBMS and the associated applications can be referred to as a "database system". Often the term "database" is also used loosely to refer to any of the DBMS, the database models that they support. Relational databases became dominant in the 1980s. These model data as rows and columns in a series of tables, and the vast majority use SQL for writing and querying data. In the 2000s, non-relational databases became popular, referred to as NoSQL because they use different query languages. Terminology and overview Formally, a "database" refers to a set of related data and the way it is organized. Access to this data is usually provided by a "database management system" (DBMS) consisting of an integrated set of computer software that allows users to interact with one or more databases and provides access to all of the data contained in the database (although restrictions may exist that limit access to particular data). The DBMS provides various functions that allow entry, storage and retrieval of large quantities of information is organized. Because of the close relationship between them, the term "database" is often used casually to refer to both a database and the DBMS used to manipulate it. Outside the world of professional information technology, the term database is often used to refer to any collection of related data (such as a spreadsheet or a card index) as size and usage requirements typically necessitate use of a database management of a database and its data which can be classified into four main functional groups: Data definition - Creation, modification and removal of definitions that define the organization of the data. Update - Insertion, modification, and deletion of the actual data.[2] Retrieval - Providing information in a form directly usable or for further processing by other applications. The retrieved data may be made available in a form basically the same as it is stored in the database or in a new form obtained by altering and monitoring users, enforcing data security, monitoring data integrity, dealing with concurrency control, and recovering information that has been corrupted by some event such as an unexpected system refers collectively to the database model.[5] "Database model.[6] Physically, database model.[6] Physi servers are dedicated computers that hold the actual databases and run only the DBMS and related software. Database servers are usually multiprocessor computers, with generous memory and RAID disk arrays used for stable storage. Hardware database accelerators, connected to one or more servers via a high-speed channel, are also used in large volume transaction processing environments. DBMSs are found at the heart of most database applications. DBMSs may be built around a custom multitasking kernel with built-in networking support, but modern DBMSs typically rely on a standard operating system to provide these functions. [citation needed] Since DBMSs comprise a significant market, computer and storage vendors often take into account DBMS requirements in their own development plans.[7] Databases and DBMSs can be categorized according to the database model(s) that they support (such as relational or XML), the type(s) of computer they run on (from a server cluster to a mobile phone), the query language(s) used to access the database (such as SQL or XQuery), and their internal engineering, which affects performance, scalability, resilience, and security. History The sizes, capabilities, and performance of databases and their respective DBMSs have grown in orders of magnitude. These performance increases were enabled by the technology progress in the areas of processors, computer memory, computer networks. The concept of a database was made possible by the emergence of direct access storage media such as magnetic tape. The subsequent development of database technology can be divided into three eras based on data model or structure: navigational,[8] SQL/relational, and post-relational. The two main early navigational data models were the hierarchical model and the CODASYL model (network model). These were characterized by the use of pointers (often physical disk addresses) to follow relationships from one record to another. The relational model, first proposed in 1970 by Edgar F. Codd, departed from this tradition by insisting that applications should search for data by content, rather than by following links. The relational model employs sets of ledger-style tables, each used for a different type of entity. Only in the mid-1980s did computing hardware become powerful enough to allow the wide deployment of relational systems (DBMSs plus applications). By the early 1990s, however, relational systems dominated in all large-scale data processing applications). By the early 1990s, however, relational systems (DBMSs plus applications) and as of 2018[update] they remain dominant: IBM DB2, Oracle, MySQL, and Microsoft SQL Server are the most searched DBMS.[9] The dominant database language, standardised SQL for the relational model, has influenced database languages for other data models.[citation needed] Object databases were developed in the 1980s to overcome the inconvenience of object-relational impedance mismatch, which led to the coining of the term "post-relational" and also the development of hybrid object-relational databases. The next generation of post-relational databases in the late 2000s became known as NoSQL databases in the late 2000s became known as NoSQL databases. relational/SQL model while aiming to match the high performance of NoSQL compared to commercially available relational DBMSs. 1960s, navigational DBMSs. 1960s, navigational DBMSs. storage (disks and drums) from the mid-1960s onwards. The term represented a contrast with the tape-based systems of the past, allowing shared interactive use rather than daily batch processing. The Oxford English Dictionary cites a 1962 report by the System Development Corporation of California as the first to use the term "data-base" in a specific technical sense.[10] As computers grew in speed and capability, a number of general-purpose database systems emerged; by the mid-1960s a number of such systems had come into commercial use. Interest in a standard began to grow, and Charles Bachman, author of one such product, the Integrated Data Store (IDS), founded the Database Task Group within CODASYL, the group responsible for the creation and standardization of COBOL. In 1971, the Database Task Group delivered their standard, which generally became known as the CODASYL approach offered applications the ability to navigate around a linked data set which was formed into a large network. Applications could find records by one of three methods: Use of a primary key (known as a CALC key, typically implemented by hashing) Navigating relationships (called sets) from one record to another Scanning all the records in a sequential order Later systems added B-trees to provide alternate access paths. Many CODASYL databases also added a declarative query language for end users (as distinct from the navigational API). However CODASYL databases were complex and required significant training and effort to produce useful applications. IBM also had their own DBMS in 1966, known as Information Management System (IMS). IMS was a development of software written for the Apollo program on the System/360. IMS was generally similar in concept to CODASYL, but used a strict hierarchy for its model of data navigation instead of CODASYL, but used a strict hierarchy for its model of data navigation instead of CODASYL is network model. way data was accessed: the term was popularized by Bachman's 1973 Turing Award presentation The Programmer as Navigator. IMS is classified by IBM as a hierarchical database are classified as network databases. IMS remains in use as of 2014[update].[11] 1970s, relational DBMS Edgar F. Codd worked at IBM in San Jose, California, in one of their offshoot offices that was primarily involved in the development of hard disk systems. He was unhappy with the navigational model of the CODASYL approach, notably the lack of a "search" facility. In 1970, he wrote a number of papers that outlined a new approach to database construction that eventually culminated in the groundbreaking A Relational Model of Data for Large Shared Data Banks.[12] In this paper, he described a new system for storing and working with large databases. Instead of records being stored in some sort of linked list of free-form records as in CODASYL, Codd's idea was to organize the data as a number of "tables" each table being used for a different type of entity. Each table would contain a fixed number of columns containing the attributes of the entity. One or more columns of each table were designated as a primary key by which the rows of the table could be uniquely identified; cross-references between tables always used these primary keys, rather than disk addresses, and queries would join tables based on these key relationships, using a set of operations based on the mathematical system of relations) aimed to ensure that each "fact" was only stored once, thus simplifying update operations. Virtual tables called views could present the data in different users, but views could not be directly updated. Codd used mathematical terms to define the model: relations, tuples, and domains rather than tables, rows, and columns. The terminology that is now familiar came from early implementations. Codd would later criticize the tendency for practical implementations to depart from the mathematical foundations on which the model was based. In the relational model, records are "linked" using virtual keys not stored in the database but defined as needed between the database but defined as needed between the data contained in the relational model, records are "linked" using virtual keys not stored in the database but defined as needed between the data contained in the relational model, records are "linked" using virtual keys not stored in the database but defined as needed between the data contained in the relational model, records are "linked" using virtual keys not stored in the database but defined as needed between the data contained in the relational model, records are "linked" using virtual keys not stored in the database but defined as needed between the data contained in the relational model, records are "linked" using virtual keys not stored in the database but defined as needed between the data contained in the relational model, records are "linked" using virtual keys not stored in the database but defined as needed between the data contained in the relational model, records are "linked" using virtual keys not stored in the database but defined as needed between the database but defined as needed b cross-table relationships, rather than disk addresses, had two primary motivations. From an engineering perspective, it enabled tables to be relocated and resized without expensive database reorganization. But Codd was more interested in the difference in semantics: the use of explicit identifiers made it easier to define update operations with clean mathematical definitions, and it also enabled query operations to be defined in terms of the established discipline of first-order predicate calculus; because these operations have clean mathematical properties, it becomes possible to rewrite queries in provably correct ways, which is the basis of query optimization. There is no loss of expressiveness compared with the hierarchic or network models, though the connections between tables are no longer so explicit. In the hierarchic and network models, records were allowed to have a complex internal structure. For example, the salary history of an employee might be represented as a "repeating group" within the employee record. In the relational model, the process of normalization led to such internal structures being replaced by data held in multiple tables, connected only by logical keys. For instance, a common use of a database system is to track information, various addresses and phone numbers. In the navigational approach, all of this data would be placed in a single variable-length record. In the relational approach, the data would be created in these optional tables only if the address or phone numbers were actually provided. As well as identifying rows/records using logical identifiers rather than disk addresses, Codd changed the way in which applications assembled data from multiple records. Rather than requiring applications to gather data one record at a time by navigating the links, they would use a declarative query language that expressed what data was required, rather than the access path by which it should be found. Finding an efficient access path to the data became the responsibility of the database management system, rather than the application, depended on the fact that queries were expressed in terms of mathematical logic. Codd's paper was picked up by two people at Berkeley, Eugene Wong and Michael Stonebraker. They started a project known as INGRES using funding that had already been allocated for a geographical database project and student programmers to produce code. Beginning in 1973, INGRES delivered its first test products which were generally ready for widespread use in 1979. INGRES was similar to System R in a number of ways, including the use of a "language" for data access, known as QUEL. Over time, INGRES moved to the emerging SQL standard. IBM itself did one test implementation of the relational model, PRTV, and a production one, Business System 12, both now discontinued. Honeywell wrote MRDS for Multics, and now there are two new implementations: Alphora Dataphor and Rel. Most other DBMS implementations usually called relational are actually SQL DBMSs. In 1970, the University of Michigan began development of the MICRO Information Management System[13] based on D.L. Childs' Set-Theoretic Data model.[14][15][16] MICRO was used to manage very large data sets by the US Department of Labor, the U.S. Environmental Protection Agency, and researchers from the University of Alberta, the University of Michigan Terminal System.[17] The system remained in production until 1998. Integrated approach Main article: Database machine In the 1970s and 1980s, attempts were made to build database systems with integrated hardware and software. The underlying philosophy was that such integrated hardware and software. The underlying philosophy was that such integrated hardware and software. to hardware support for database management was ICL's CAFS accelerator, a hardware disk controller with programmable search capabilities. In the long term, these efforts were generally unsuccessful because specialized database machines could not keep pace with the rapid development and progress of general-purpose computers. Thus most database systems nowadays are software systems running on general-purpose hardware, using general-purpose computer data storage. However, this idea is still pursued for certain applications by some companies like Netezza and Oracle (Exadata). Late 1970s, SQL DBMS IBM started working on a prototype system loosely based on Codd's concepts as System R in the early 1970s. The first version was ready in 1974/5, and work then started on multi-table systems in which the data could be split so that all of the data c a standardized query language - SQL[citation needed] - had been added. Codd's ideas were establishing themselves as both workable and superior to CODASYL, pushing IBM to develop a true production version of System R, known as SQL/DS, and, later, Database 2 (DB2). Larry Ellison's Oracle Database (or more simply, Oracle) started from a different chain, based on IBM's papers on System R. Though Oracle V1 implementations were completed in 1978, it wasn't until Oracle Version 2 when Ellison beat IBM to market in 1979.[18] Stonebraker went on to apply the lessons from INGRES to develop a new database, Postgres, which is now known as PostgreSQL. PostgreSQL is often used for global mission-critical applications (the .org and .info domain name registries use it as their primary data store, as do many large companies and financial institutions). In Sweden, Codd's paper was also read and Mimer SQL was developed from the mid-1970s at Uppsala University. In 1984, this project was consolidated into an independent enterprise. Another data model, the entity-relationship model, and the difference between the two as a data modeling constructs were retrofitted as a data modeling construct for the relational model, and the difference between the two have become irrelevant.[citation needed] 1980s, on the desktop The 1980s ushered in the age of desktop computing. The new computers empowered their users with spreadsheets like Lotus 1-2-3 and database software like dBASE. The dBASE product was lightweight and easy for any computer user to understand out of the box. C. Wayne Ratliff, the creator of dBASE, stated: "dBASE was different from programs like BASIC, C, FORTRAN, and COBOL in that a lot of the dirty work had already been done. The data manipulation is done by dBASE instead of by the user, so the user can concentrate on what he is doing, rather than having to mess with the dirty details of opening, reading, and closing files, and managing space allocation."[19] dBASE was one of the top selling software titles in the 1980s, along with a rise in object-oriented programming, saw a growth in how data in various databases were handled. Programmers and designers began to treat the data in their databases as objects. That is to say that if a person's data were in a database, that person's attributes, such as their address, phone number, and age, were now considered to belong to that person instead of being extraneous data. This allows for relational fields.[20] The term "object-relational fields.[2 impedance mismatch" described the inconvenience of translating between programmed objects and databases attempt to solve this problem by providing an object-oriented language (sometimes as extensions to SQL) that programmers can use as alternative to purely relational SQL. On the programming side, libraries known as object-relational mappings (ORMs) attempt to solve the same problem. 2000s, NoSQL and NewSQL XML databases are a type of structured document-oriented database that allows querying based on XML document attributes. XML databases are mostly used in applications where the data is conveniently viewed as a collection of documents, with a structure that can vary from the very flexible to the highly rigid: examples include scientific articles, patents, tax filings, and personnel records. NoSQL databases are often very flexible to the highly rigid: examples include scientific articles, patents, tax filings, and personnel records. are designed to scale horizontally. In recent years, there has been a strong demand for massively distributed system to simultaneously provide consistency, availability, and partition tolerance guarantees. A distributed system can satisfy any two of these guarantees at the same time, but not all three. For that reason, many NoSQL databases are using what is called eventual consistency to provide both availability and partition tolerance guarantees with a reduced level of data consistency. NewSQL is a class of modern relational databases that aims to provide the same scalable performance of NoSQL systems for online transaction processing (read-write) workloads while still using SQL and maintaining the ACID guarantees of a traditional database system. Use cases This section does not cite any sources. Please help improve this section does not cite any sources. 2013) (Learn how and when to remove this template message) Databases are used to support internal operations of organizations and to underpin online interactions with customers and suppliers (see Enterprise software). Databases are used to hold administrative information and more specialized data, such as engineering data or economic models (see Enterprise software). Examples include computerized library systems, flight reservation systems, computerized parts inventory systems, and many content management systems, and many content management systems, and many content management systems. or multimedia objects. Another way is by their application area, for example: accounting, music compositions, movies, banking, manufacturing, or insurance. A third way is by some technical aspect, such as the databases. An in-memory database is a database that primarily resides in main memory, but is typically backed-up by non-volatile computer data storage. Main memory databases, and so are often used where response time is critical, such as in telecommunications network equipment. An active databases, and so are often used where response time is critical, such as in telecommunications network equipment. which can respond to conditions both inside and outside the database. Possible uses include security monitoring, alerting, statistics gathering and authorization. Many database features in the form of database triggers. A cloud database relies on cloud technology. Both the database and most of its DBMS reside remotely, "in the cloud", while its applications are both developed by programmers and later maintained and used by end-users through a web browser and Open APIs. Data warehouses archive data for use by managers and other end-users who may not have access to operational data. For example, sales data might be aggregated to weekly totals and converted from internal product codes to use UPCs so that they can be compared with ACNielsen data. transforming, loading, and managing data so as to make them available for further use. A deductive database combines logic programming with a relational database is one in which both the data and the DBMS span multiple computers. A document-oriented database is designed for storing, retrieving, and managing documentoriented, or semi structured, information. Document-oriented databases are one of the main categories of NoSQL databases. An embedded database system is a DBMS which is tightly integrated with an application's end-users and requires little or no ongoing maintenance.[21] End-user databases consist of data developed by individual end-users. Examples of these are collections of documents, spreadsheets, presentations, multimedia, and other files. Several products exist to support such databases. functionality. A federated database system comprises several distinct databases, each with its own DBMS, possibly of different types (in which case it would also be a heterogeneous database system), which transparently integrates multiple autonomous DBMSs, possibly of different types (in which case it would also be a heterogeneous database system). and provides them with an integrated conceptual view. Sometimes the term multi-database is used as a synonym to federated database, though it may refer to a less integrated (e.g., without an FDBMS and a managed integrated schema) group of databases that cooperate in a single application. In this case, typically middleware is used for distribution, which typically includes an atomic commit protocol (ACP), e.g., the two-phase commit protocol, to allow distributed (global) transactions across the participating databases. A graph database that uses graph structures with nodes, edges, and properties to represent and store information. General graph databases that can store any graph are distinct from specialized graph databases. An array DBMS is a kind of NoSQL DBMS that allows modeling, storage, and retrieval of (usually large) multi-dimensional arrays such as satellite images and climate simulation output. In a hypertext or hypermedia database any word or a piece of text representing an object, e.g., another piece of text, an article, a picture, or a film, can be hyperlinked to that object. Hypertext databases are particularly useful for organizing large amounts of disparate information. For example, they are useful for organizing online encyclopedias, where users can conveniently jump around the text. The World Wide Web is thus a large distributed hypertext database. A knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge management, providing the means for the computerized collection, organization, and retrieval of knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database for knowledge base (abbreviated KB, kb or Δ[22][23]) is a special kind of database (abbreviated KB, kb or Δ[22][23]) is a special kind of database (abbreviated KB, kb or Δ[22][23]) is a special kind of da and related experiences. A mobile database can be carried on or synchronized from a mobile computing device. Operational databases store detailed data about the operations. Examples include customer databases that record contact, credit, and demographic information about a business's customers, personnel databases that hold information such as salary, benefits, skills data about product components, parts inventory, and financial databases that keep track of the organization's money, accounting and financial dealings. A parallel database seeks to improve performance through parallelization for tasks such as loading data, building indexes and evaluating queries. The major parallel DBMS architectures which are induced by the underlying hardware architectures which are induced by the underlying hardware architecture are: well as other data storage. Shared disk architecture, where each processing unit has its own main memory, but all units share the other storage. Shared-nothing architecture, where each processing unit has its own main memory, but all units share the other storage. inferences from imprecise data. Real-time databases process transactions fast enough for the result to come back and be acted on right away. A spatial database can store the data with multidimensional features. The queries on such data include location-based queries, like "Where is the closest hotel in my area?". A temporal database has built-in time aspects, for example a temporal data model and a temporal version of SQL. More specifically the temporal aspects usually include valid-time and transaction-time. A terminology-oriented database builds upon an object-oriented database builds upon an object-oriented database. protected way diverse objects that do not fit naturally and conveniently in common databases. It may include email messages, documents, journals, multimedia objects can be highly structured. However, the entire possible object collection does not fit into a predefined structured framework. Most established DBMSs now support unstructured data in various ways, and new dedicated DBMSs are emerging. Database management system (DBMS) as a "software system that enables users to define, create, maintain and control access to the database".[24] Examples of DBMS's include MySQL, PostgreSQL, Microsoft SQL Server, Oracle Database, and Microsoft Access. The DBMS for the object (oriented) and ORDBMS for the object (oriented) and ORDBMS for the object and or characteristic, such as DDBMS for a distributed database management systems. The functionality provided by a DBMS can vary enormously. The core functionality is the storage, retrieval and update User accessible catalog or data dictionary describing the metadata Support for transactions and concurrency Facilities for recovering the database should it become damaged Support for authorization of access and update of data Access support for transactions and concurrency Facilities for recovering the database should it become damaged Support for authorization of access and update of data Access support for transactions and concurrency Facilities for recovering the database should it become damaged Support for transactions and concurrency Facilities for recovering the database should it become damaged Support for transactions and concurrency Facilities for recovering the database should it become damaged Support for transactions and concurrency Facilities for recovering the database should it become damaged Support for transactions and concurrency Facilities for recovering the database should it become damaged Support for transactions and concurrency Facilities for recovering the database should it become damaged Support for transactions and concurrency Facilities for recovering the database should it become damaged Support for transactions and concurrency Facilities for recovering the database should it become damaged Support for transactions and concurrency Facilities for recovering the database should be access and update of data Access support for transactions and concurrency facilities for recovering the database should be access and update of data Access support for transactions and concurrency facilities for transactions and concurrency facili be expected the DBMS will provide a set of utilities for such purposes as may be necessary to administer the database effectively, including import, export, monitoring, defragmentation and analysis utilities. [26] The core part of the DBMS interacting between the database engine Often DBMSs will have configuration parameters that can be statically and dynamically tuned, for example the maximum amount of manual configuration, and for cases such as embedded databases the need to target zero-administration is paramount. The large major enterprise DBMSs have tended to increase in size and functionality and can have involved thousands of human years of development effort through their lifetime.[a] Early multi-user DBMS typically only allowed for the application to reside on the same computer with access via terminals or terminal emulation software. The client-server architecture was a development where the application resided on a client desktop and the database on a server allowing the processing to be distributed. This evolved into a multitier architecture incorporating application servers and web servers with the end user interface via a web browser with the database only directly connected to the adjacent tier.[27] A general-purpose DBMS will provide public application programming interfaces (API) and optionally a processor for database. A special purpose DBMS may use a private API and be specifically customized and linked to a single application. For example, an email system performing many of the functions of a general-purpose DBMS such as message insertion, message deletion, attachment handling, blocklist lookup, associating messages an email address and so forth however these functions are limited to what is required to handle email. Application Main article: Database application External interaction with the database will be via an application program that interfaces with the DBMS.[28] This can range from a database tool that allows users to execute SQL queries textually or graphically, to a web site that happens to use a database tool that allows users to execute SQL queries textually or graphically. the database (sometimes referred to as a datasource) via an application program interface (API) or via a database language. The particular API or language chosen will need to be supported by DBMS, possible indirectly via a preprocessor or a bridging API. Some API's aim to be database independent, ODBC being a commonly known example. Other common API's include JDBC and ADO.NET. Database languages are special-purpose languages; Data control languages, which allow one or more of the following tasks, sometimes distinguished as sublanguages; Data control languages are special-purpose languages. dropping tables and the relationships among them; Data manipulation language (DML) - performs tasks such as inserting, updating, or deleting data occurrences; Data query language are specific to a particular data model. Notable examples includes SQL combines the roles of data definition, data manipulation, and query in a single language. It was one of the first commercial languages for the relational model, although it departs in some respects from the relational model as described by Codd (for example, the rows and columns of a table can be ordered). SQL became a standard of the American National Standards Institute (ANSI) in 1986, and of the International Organization (ISO) in 1987. The standards have been regularly enhanced since and is supported (with varying degrees of conformance) by all mainstream commercial relational DBMSs.[29][30] OQL is an object model language standard (from the Object Data Management Group). It has influenced the design of some of the newer query languages like JDOQL and EJB QL. XQuery is a standard XML query language implemented by XML database systems such as MarkLogic and eXist, by relational databases with XML capability such as Oracle and DB2, and also by in-memory XML processors such as Saxon. SQL/XML combines XQuery with SQL.[31] A database language may also incorporate features like: DBMS-specific configuration and storage engine management Computations to modify query results, like counting, summing, averaging, sorting, grouping, and cross-referencing Constraint enforcement (e.g. in an automotive database only allowing one engine type per car) Application programming interface version of the query language, for programmer convenience Storage is the container of the physical materialization of a database. It comprises the internal (physical) level in the database architecture. It also contains all the information needed (e.g., metadata, "data about the data", and internal level when needed. Putting data into permanent storage is generally the responsibility of the database engine a.k.a. "storage engine". Though typically accessed by a DBMS through the underlying operating system (and often using the operating systems' file systems as intermediates for storage layout), storage properties and configuration setting are extremely important for the efficient operation of the DBMS, and thus are closely maintained by database administrators. A DBMS, while in operation always has its database residing in several types of storage (e.g., memory and external storage). The database data and the additional needed information, possibly in very large amounts, are coded into bits. Data typically reside in the storage in structures that look completely different from the way the data look in the conceptual and external levels, but in ways that attempt to optimize (the best possible) these levels' reconstruction when needed by users and programs, as well as for computing additional types of needed information from the data (e.g., when querying the database). Some DBMSs support specifying which character encoding was used to store data, so multiple encodings can be used in the same database. Various low-level database storage structures are used by the storage engine to serialize the data model so it can be written to the medium of choice. Techniques such as indexing may be used to improve performance. Conventional storage is row-oriented, but there are also column-oriented and correlation databases Materialized views Main article: Materialized views, which consist of frequently needed external views or query results. Storing such views saves the expensive computing of them each time they are needed. The downsides of materialized views are the overhead incurred when updating them to keep them synchronized with their original updated database employs storage redundancy. Replication Occasionally a database employs storage redundancy. Replication Main article: Database data availability (both to improve performance of simultaneous multiple end-user accesses to a same database object, and to provide resiliency in a case of partial failure of a distributed database). Updates of a replicated object need to be synchronized across the object copies. In many cases, the entire database object, and to provide resiliency in a case of partial failure of a distributed database). the article Database security. Please see discussion on the linked talk page. (March 2013) (Learn how and when to remove this template message) Main article: Database security deals with all various aspects of protecting the database security deals with all various aspects of protecting the database security. database uses to unintentional database access by unauthorized entities (e.g., a person or a computer program). Database access what information in the database. The information may comprise specific database objects (e.g., record types, specific database). records, data structures), certain computations over certain objects (e.g., query types, or specific queries), or using specific access paths to the former (e.g., using specific access paths to the former (e.g., using specific access paths to the former (e.g., using specific access paths), or using specific access paths to the former (e.g., using specific access paths), or using specific access paths to the former (e.g., using specific access paths), or using specific access paths to the former (e.g., using specific access paths), or using specific access paths to the former (e.g., using specific access paths), or using specific access paths to the former (e.g., using specific access paths), or using specific access paths to the former (e.g., using specific access paths), or using specific access paths to the former (e.g., using specific access paths), or using specific access paths to the former (e.g., using specific access paths), or using specific access paths), or using specific access paths), or using specific access paths to the former (e.g., using specific access), or using specific access), or using specific access paths), or using specific access), or usi security DBMS interfaces. This may be managed directly on an individual basis, or by the assignment of individuals and groups to roles which are then granted entitlements. Data security prevents unauthorized users from viewing or updating the database. Using passwords, users are allowed access to the entire database or subsets of it called "subschemas". For example, an employee, but one group of users may be authorized to view only payroll data, while others are allowed access to only work history and medical data. If the DBMS provides a way to interactively enter and update the database, as well as interrogate it, this capability allows for managing personal databases. Data security), or the interpretation of them, or parts of them to meaningful information (e.g., by looking at the strings of bits that they comprise, concluding specific valid credit-card numbers; e.g., see data encryption). Change and access logging records who accessed which attributes, what was changed, and when it was changed at a encryption). record of access occurrences and changes. Sometimes application-level code is used to record changes rather than leaving this to the database. Monitoring can be set up to attempt to detect security breaches. Transactions and concurrency Further information: Concurrency control Database transactions can be used to introduce some level of fault tolerance and data integrity after recovery from a crash. A database transaction is a unit of work, typically encapsulating a number of operations over a database and also other systems. Each transaction has well defined boundaries in terms of which program/code executions are included in that transaction (determined by the transaction's programmer via special transaction commands). The acronym ACID describes some ideal properties of a database transaction; atomicity, consistency, isolation, and durability. Migration See also: Data migration § Database migration A database built with one DBMS is not portable to another DBMS (i.e., the other DBMS cannot run it). However, in some situations, it is desirable to migrate a database from one DBMSs may have different total costs of ownership or TCOs), functional, and operational (different DBMSs may have different total costs). capabilities). The migration involves the database related application from one DBMS type to another. The transformation from one DBMS type to another. The transformation from one DBMS type to another. desired that also some aspects of the architecture internal level are maintained. A complex or large database migration to migrate. This in spite of the fact that tools may exist to help migration between specific DBMSs. Typically, a DBMS vendor provides tools to help importing databases from other popular DBMSs. Building, maintaining, and tuning Main article: Database tuning After designing a database tor an appropriate general-purpose. A DBMS provides the needed user interfaces to be used by database administrators to define the needed application's data structures within the DBMS's respective data model. Other user interfaces are used to select needed components (like security related, storage allocation parameters, etc.). When the database is ready (all its data structures and other needed components) and other needed components (like security related, storage allocation parameters, etc.). are defined), it is typically populated with initial application's data (database initialization, which is typically a distinct project; in many cases using specialized DBMS interfaces that support bulk insertion) before making it operational. In some cases, the database becomes operational while empty of application data, and data are accumulated during its operation. After the database is created, initialized and populated it needs to be maintained. Various database parameters may need to be tuned (tuning) for better performance; application's data structures may be written to add to the application's data structures may be changed or added, new related application programs may be written to add to the application's data structures may be changed or added, new related application programs may be written to add to the application's data structures may be changed or added, new related application's data structures may be changed or added, new related application's data structures may be changed or added, new related application's data structures may be written to add to the application's data structures may be changed or added, new related application's data structures may be written to add to the applic functionality, etc. Backup and restore Main article: Backup Sometimes it is desired to bring a database back to a previous state (for many reasons, e.g., cases when the database is found corrupted due to a software error, or if it has been updated with erroneous data). To achieve this, a backup operation is done occasionally or continuously, where each desired database state (i.e., the values of its data and their embedding in database's data structures) is kept within dedicated backup files (many techniques exist to do this state by a desired point in time when the database's data structures) is kept within dedicated backup files (many techniques exist to do this effectively). When it is decided by a database back to this state by a desired point in time when the database's data structures) is kept within dedicated backup files (many techniques exist to do this effectively). was in this state), these files are used to restore that state. Static analysis Static analysis techniques for software verification can be applied also in the scenario of query languages. In particular, the *Abstract interpretation framework has been extended to the field of query languages for relational databases as a way to support sound approximation techniques.[32] The semantics of query languages can be tuned according to suitable abstractions of the concrete domain of data. The abstractions, in particular, for security purposes, such as fine grained access control, watermarking, etc. Miscellaneous features Other DBMS features might include: Database logs - This helps in keeping a history of the executed functions. Graphics component for producing graphs and charts, especially in a data warehouse system. Query optimization on every query to choose an efficient query plan (a partial order (tree) of operations) to be executed to compute the query result. May be specific to a particular storage engine. Tools or hooks for database design, application programming, application programming units) and related database mapping (especially for a distributed DBMS), storage allocation and database layout monitoring, storage migration, etc. Increasingly, there are calls for a single system that incorporates all of these core functionalities into the same build, test, and deployment framework for database management and source control. Borrowing from other developments in the software industry, some market such offerings as "DevOps for database".[33] Design and modeling Main article: Database designer is to produce a conceptual data model that reflects the structure of the information to be held in the database. A common approach to this is to develop an entity-relationship model, often with the aid of drawing tools. Another popular approach is the Unified Modeling Language. A successful data model will accurately reflect the possible state of the external world being modeled: for example, if people can have more than one phone number, it will allow this information to be captured. Designing a good conceptual data model requires a good understanding of the application, like "can a customer also be a supplier?", or "if a product is sold with two different forms of packaging, are those the same products?", or "if a product set of a customer also be a supplier?", or "if a product is sold with two different forms of packaging, are those the same product set of a customer also be a supplier?", or "if a product is sold with two different forms of packaging, are those the same product or different forms of packaging, are those the same product set of a customer also be a supplier?", or "if a product is sold with two different forms of packaging, are those the same product or different forms of packaging, are those the same product is sold with two different forms of packaging, are those the same product or different forms of packaging, are those the same product or different forms of packaging, are those the same product or different forms of packaging, are those the same product or different forms of packaging, are those the same product or different forms of packaging, are those the same product or different forms of packaging, are those the same product or different forms of packaging, are those the same product or different forms of packaging, are those the same product or different forms of packaging, are those the same product or different forms of packaging, are those the same product or different forms of packaging, are those the same product of the same packaging are those t plane flies from New York to Dubai via Frankfurt, is that one flight or two (or maybe even three)?". The answers to these questions establish definitions of the terminology used for entities (customers, products, flights, flight segments) and their relationships and attributes. processes, or the analysis of workflow in the organization. This can help to establish what information is needed in the database, and what can be left out. For example, it can help when deciding whether the database needs to hold historic data as well as current data. is to translate this into a schema that implements the relevant data structures within the database. This process is often called logical data model is (in theory at least) independent of the choice of database technology, the logical data model will be expressed in terms of a particular database model are often used to express that design). The most popular database model for the design of a specific database, and database model for the design of a specific database model are often used to express that design). model for general-purpose databases is the relational model, or more precisely, the relational model as represented by the SQL language. The process of creating a logical database design using this model uses a methodical approach known as normalization. The goal of normalization is to ensure that each elementary "fact" is only recorded in one place, so that insertions, updates, and deletions automatically maintain consistency. The final stage of database design is to make the decisions that affect performance, scalability, recovery, security, and the like, which depend on the particular DBMS. This is often called physical database design, and the output is the physical data model. A key goal during this stage is data independence, meaning that the decisions made for performance optimization purposes should be invisible to end-users and applications. There are two types of data independence: Physical data independence and logical data independence and logical data independence. knowledge of the expected workload and access patterns, and a deep understanding of the features offered by the chosen DBMS. Another aspect of physical database model as defining security levels and methods for the data itself. Models Main article: Database model Collage of five types of database model is a type of data model that determines the logical structure of a database and fundamentally determines in which manner data can be stored, organized, and manipulated. The most popular example of a database model is the relational model (or the SQL approximation of relational), which uses a table-based format. Common logical data model Entity-relationship model Entity-relationship model Entity-relationship model Entity-relationship model Entity-relationship model Star schema An object-relational database combines the two related structures. Physical data models include: Inverted index Flat file Other model Specialized models are optimized for particular types of data: XML database Semantic model Content store Event store Time series model External, conceptual, and internal views Traditional view of data[34] A database management system provides three views of the database data: The external level defines how each group of end-users sees the organization of data in the database data: The external level defines how each group of end-users sees the organization of data in the database data: The external level defines how each group of end-users sees the organization of data in the database data: The external level defines how each group of end-users sees the organization of data in the database data: The external level defines how each group of end-users sees the organization of data in the database data: The external level defines how each group of end-users sees the organization of data in the database data. compatible global view.[35] It provides the synthesis of all the external views. It is out of the scope of the various database administrators. The internal level (or physical level) is the internal organization of data inside a DBMS. It is concerned with cost, performance, scalability and other operational matters. It deals with storage layout of the data, using storage structures such as indexes to enhance performance. Occasionally it stores data of individual views (materialized views), computed from generic data, if performance performance is storage structures such as indexes to enhance performance. requirements, possibly conflicting, in an attempt to optimize overall performance across all activities. While there is typically only one conceptual (or internal) views. This allows users to see database information in a more business-related way rather than from a technical, processing viewpoint. For example, a financial department of a company needs the payment details of all employees that are the interest of the human resources department. Thus different departments need different views of the company's database. The threelevel database architecture relates to the concept of data independence which was one of the major initial driving forces of the relational model. The idea is that changes in the internal level do not affect the view at a higher level. For example, changes in the internal level do not affect application programs written using conceptual level interfaces, which reduces the impact of making physical changes to improve performance. The conceptual view provides a level of indirection between internal and external. On one hand it provides a common view of the data are stored or managed (internal level). In principle every level, and even every external view, can be presented by a different data model. In practice usually a given DBMS uses the same data model for both the external and the conceptual levels (e.g., relational model). The internal level, which is hidden inside the DBMS and depends on its implementation requires a different level of detail and uses its own types of data structure types. Separating the external, conceptual and internal levels was a major feature of the relational database model implementations that dominate 21st century databases.[35] Research Database technology has been an active research topic since the 1960s, both in academia and in the research and development groups of companies (for example IBM Research). Research activity includes theory and development of prototypes. Notable research topics have included models, the atomic transaction concept, and related concurrency control techniques, query languages and query optimization methods, RAID, and more. The database research area has several dedicated academic journals (for example, ACM Transactions on Database Systems-TODS, Data and Knowledge Engineering-DKE) and annual conferences (e.g., ACM SIGMOD, ACM PODS, VLDB, IEEE ICDE). See also For a topical guide to this subject, see Outline of databases. Comparison of database tools Comparison of object database management systems Comparison of relational database management systems Comparison of relational database theory Database testing dataset Notes ^ This article quotes a development time of 5 years involving 750 people for DB2 release 9 alone. (Chong et al. 2007) References ^ Ullman & Widom 1997, p. 1. ^ "Update - Definition of update by Merriam-Webster". merriam-webster.com. ^ "Retrieval - Definition of retrieval by Merriam-Webster". merriam-webster.com. ^ "Administration - Definition of administration by Merriam-Webster". merriam-webster.com. ^ Tsitchizris & Lochovsky 1982. ^ Beynon-Davies 2003. ^ Nelson & Nelson 2001. ^ Bachman 1973. ^ "TOPDB Top Database index". pypl.github.io. ^ "database, n". OED Online. Oxford University Press. June 2013. Retrieved July 12, 2013. (Subscription required.) ^ IBM Corporation (October 2013). 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